

Noncomputable Coding, Density, and Stochasticity

Justin Miller

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<https://sites.nd.edu/justin-d-miller/talks/>

Asymptotic Density

Recall that the (asymptotic) density of X is

$$\rho(X) = \lim_{n \rightarrow \infty} \frac{|X \upharpoonright n|}{n}$$

if this limit exists. $\bar{\rho}(X)$ and $\underline{\rho}(X)$, the upper and lower density, are the limit superior and limit inferior respectively.

Intrinsic Density

Intrinsic Density is the computable permutation invariant form of asymptotic density, that is A has intrinsic density α if $\rho(\pi(A)) = \alpha$ for all computable permutations π . We denote it by $P(A)$ if it exists. A is intrinsically small if $P(A) = 0$.

The absolute upper density, $\overline{P}(A)$, is the limit superior of $\rho(\pi(A))$ over all computable permutations π . The absolute lower density, $\underline{P}(A)$, is the limit inferior.

Randomness

What values of $r \in [0, 1]$ are achieved as the intrinsic density of a set?

The short answer is everything: 0 and 1 are well known. If μ_r is the Bernoulli measure with parameter r (i.e. the r -biased coin flip measure), then standard arguments show that every μ_r -1-random set has intrinsic density r . (It also holds for Schnorr randomness, but the proof takes more work.)

Motivation

Can we use sets of a given intrinsic density to understand those of a different one?

As we shall see, sets with intrinsic density r are a strict super set of the μ_r -randoms and have fundamentally different closure properties as a class. (The join preserves intrinsic density.) We shall introduce noncomputable coding methods, the `into` and `within` operations, to change intrinsic densities and construct an example for every real $r \in (0, 1)$ without appealing to μ_r -randomness.

Changing Density via Noncomputable Coding

To understand intrinsic density, we set out to find some method of combining sets A and B such that the intrinsic density of the resulting set is some function of the intrinsic densities of A and B . Classical operations, notably the join, do not work. Therefore we turn to the into operation.

The Into Operation

Definition

Given two sets

$$A = \{a_0 < a_1 < a_2 < \dots\}$$

and

$$B = \{b_0 < b_1 < b_2 < \dots\}$$

we define the set $B \triangleright A$, or “ B into A ,” to be

$$\{a_{b_0} < a_{b_1} < a_{b_2} < \dots\}$$

This will be a fundamental tool for changing densities.

Examples

- If E is the set of evens, then $E \triangleright E$ is the set of multiples of 4: $e_n = 2n$, so

$$e_{e_n} = e_{2n} = 2(2n) = 4n$$

- If O is the set of odds, then $E \triangleright O$ is the set of naturals congruent to 1 mod 4:
 $o_n = 2n + 1$ and $e_n = 2n$, so

$$o_{e_n} = o_{2n} = 2(2n) + 1 = 4n + 1$$

- For any A and B ,

$$A \oplus B = (A \triangleright E) \sqcup (B \triangleright O)$$

Basic Properties

- $A = A \triangleright \omega$
- $A = \omega \triangleright A$
- $B \triangleright A \subseteq A$
- $(B \triangleright A) \sqcup (\bar{B} \triangleright A) = A$
- \triangleright is associative

Into and Asymptotic Density

Lemma

- $\bar{\rho}(B \triangleright A) \leq \bar{\rho}(A)\bar{\rho}(B)$
- $\underline{\rho}(B \triangleright A) \geq \underline{\rho}(A)\underline{\rho}(B)$

Corollary

$$\rho(B \triangleright A) = \rho(B)\rho(A)$$

Intro and Intrinsic Density

For a set X , P_X represents intrinsic density relative to X , i.e. invariance under all X -computable permutations.

Theorem

If $P(A) = \alpha$ and $P_A(B) = \beta$, then $P(B \triangleright A) = \alpha\beta$.

Powers of Two

Lemma

There is a countable, disjoint sequence of sets $\{A_i\}_{i \in \omega}$ such that $P(A_i) = \frac{1}{2^{i+1}}$.
Furthermore, $\lim_{n \rightarrow \infty} \overline{P}(\bigsqcup_{i > n} A_i) = 0$.

Proof Sketch.

Let X be 1-Random. Then by the general form of Van Lambalgen's theorem and the fact that 1-Randoms have intrinsic density $\frac{1}{2}$, the columns $X^{[n]}$ of X give us countably many sets all with intrinsic density relative to the rest. Then define $B_0 = \omega$, $A_n = \overline{X^{[n]}} \triangleright B_n$, and $B_{n+1} = X^{[n]} \triangleright B_n$. □

Arbitrary Intrinsic Density

Let $r \in (0, 1)$ and $\{A_i\}_{i \in \omega}$ be as in the previous lemma. We identify B_r with the set of bits which are 1 in the binary expansion of r . Then $\bigcup_{i \in B_r} A_i$ will have intrinsic density $\sum_{i \in B_r} \frac{1}{2^{i+1}} = r$.

Corollary

If $r \in (0, 1)$ computes a 1-Random, then r computes a set of intrinsic density r .

MWC Stochasticity

Informally, we say a set X is von Mises-Wald-Church stochastic, or MWC stochastic, for r if every infinite subsequence selected from it by a process which can only use information about whether or not $i \in X$ for $i < n$ to decide whether or not to include n in the subsequence has asymptotic density r .

Intro and Stochasticity

Lemma

If A is MWC-stochastic for α relative to B and B is MWC-stochastic for β relative to A , then $B \triangleright A$ is MWC-stochastic for $\alpha\beta$.

Intro and Randomness

Question

If A is μ_α -random relative to B and B is μ_β -random relative to A , is $B \triangleright A$ $\mu_{\alpha\beta}$ -random?

The Within Operation

The into operation has a natural dual.

Definition

Given two sets

$$A = \{a_0 < a_1 < a_2 < \dots\}$$

and

$$B = \{b_0 < b_1 < b_2 < \dots\}$$

we define the set $B \triangleleft A$, or “ B within A ”, to be $\{n : a_n \in B\}$. In other words, $B \cap A$ is some subset of A , so there is some X such that $X \triangleright A = B \cap A$. In this case, $B \triangleleft A = X$.

Example

- Let T be the multiple of three. Then $T \triangleleft E = T$:

$$T \triangleleft E = \{n : 2n \in T\} = \{n : 6 \mid 2n\} = \{n : 3 \mid n\}$$

- $T \triangleleft O$ is the set of naturals congruent to 1 mod 3:

$$T \triangleleft O = \{n : 2n + 1 \in T\} = \{n : 2n \equiv 2 \pmod{3}\} = \{n : n \equiv 1 \pmod{3}\}$$

Within: Basic Properties

- $\omega = A \triangleleft A$
- $(B \triangleleft A) \sqcup (\bar{B} \triangleleft A) = \omega$
- If $B \subseteq A$, then $(B \triangleleft A) \triangleright A = B \cap A = B$.
- \triangleleft is not associative: Let E be the evens, O the odds, and N the set of naturals congruent to 2 mod 4. Then

$$(O \triangleleft N) \triangleleft E = \emptyset \triangleleft E = \emptyset$$

but

$$O \triangleleft (N \triangleleft E) = O \triangleleft O = \omega$$

Applications of Within

Let C be a computable set.

Lemma

If A has intrinsic density α , then so does $A \triangleleft C$.

Lemma

If A has MWC-stochasticity for α , then so does $A \triangleleft C$.

Lemma

If A is μ_α -random, then so is $A \triangleleft C$.