Nonexistence of Nash Equilibria in Games Over Finitely Additive Measure Spaces

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Background and Motivation

- Games with finite number of players.
- Modeling individual negligibility: Atomless measures, infinitesimals, Loeb spaces, Finitely–additive measures.
 - Standard model: continuum of players with atomless distribution.
 - Countably many agents.
 - Each player has zero mass. Measure of the whole space is 1. So, finitely additive measures.

- Nash: Every finite-player, finite-action game has an equilibrium in mixed strategies.
- Games with infinitely many players, compact convex action set of each player, payoffs are quasi-concave in own argument. There is a Nash equilibrium.
 Ma (1969)
- Atomless, countably additive measure space of players:
 - If the set of players is an atomless, countably additive measure space then a game has a pure strategy Nash equilibrium.

Schmeidler (1973)

- DWW (1951) theorem: Every mixed strategy Nash equilibrium can be purified.
- ► The DWW theorem holds for finitely additive measure spaces.
- Question: Does a pure/mixed strategy Nash equilibrium exist in a game over a finitely additive measure space of players?

- If the set of players is endowed with a finitely additive measure, then:
 - a game may not have a Nash equilibrium (in pure or mixed strategies).
 - the Nash equilibrium correspondence may not have the closed graph property.
 - a game may not have an ϵ -equilibrium.
 - sufficient conditions for existence of an ϵ -equilibrium.

Primary reason for nonexistence of equilibrium: The lack of upper hemicontinuity of the integral of a correspondence.

Large Games

- ► Let $E = \{e^1, \dots, e^L\}$ be the set of unit vectors in \mathbb{R}^L and $S = \{s \in \mathbb{R}_+^L : \sum_{k=1}^L s_k = 1\}$ the unit simplex in \mathbb{R}^L .
- ► Let U be the set of real valued continuous functions defined on E × S, endowed with sup norm.
- Let (T, T, μ) be an atomless, countably additive probability space.
- ► A (non-anonymous large) game is a measurable function $\mathcal{G} : T \longrightarrow \mathcal{U}.$
- A *pure strategy profile* (of a game \mathcal{G}) is a measurable function $f: T \longrightarrow E$.
- A f : T → E is a (pure strategy) Nash equilibrium of G if for almost all t,

$$\mathcal{G}(t)\left(f(t),\int_{\mathcal{T}}f\;\mathsf{d}\mu\right)\geq\mathcal{G}(t)\left(\mathsf{a},\int_{\mathcal{T}}f\;\mathsf{d}\mu\right)\;\text{for all}\;\mathsf{a}\in\mathsf{E}.$$

Existence of Nash Equilibrium

Theorem (Schmeidler)

Every game has a pure strategy Nash equilibrium.

• Define a correspondence $B : T \times S \longrightarrow E$ by

 $B(t,s) = \{e^k \in E: \ \mathcal{G}(t)(e^k,s) \geq \mathcal{G}(t)(a,s) \ \text{ for all } a \in E\}.$

- ▶ B(t,s) is nonempty, $B(\cdot,s)$ is measurable and $B(t,\cdot)$ is uhc.
- Let $\Gamma(s) = \int_T B(\cdot, s) d\mu$.
 - $\Gamma(s)$ is nonempty for each $s \in S$.
 - $\Gamma(\cdot)$ is uhc (integration preserves uhc).
 - $\Gamma(\cdot)$ is convex valued (by Lyapunov's theorem).
- ▶ **Γ** has a fixed point *s*^{*} (by Kakutani's fixed point theorem).
- ▶ So, there is $f : T \longrightarrow E$ such that $\int_T f d\mu = s^*$ and for almost all $t, f(t) \in B(t, s^*)$.
- ▶ This *f* is a Nash equilibrium of *G*.

Finitely Additive Measures

- ▶ *T* is a nonempty set and *T* a field of subsets of *T*. (*i*) Ø, *T* ∈ *T*; (*ii*) *A*, *B* ∈ *T* ⇒ *A* ∪ *B* ∈ *T* and (*iii*) *A*, *B* ∈ *T* ⇒ *A* \ *B* ∈ *T*.
- ▶ μ is a finitely additive probability measure on \mathcal{T} if (i) $\mu(\emptyset) = 0$, $\mu(\mathcal{T}) = 1$, $\mu(A) \ge 0$ for all $A \in \mathcal{T}$ and (ii) $\mu(A \cup B) = \mu(A) + \mu(B)$ if $A, B \in \mathcal{T}$, $A \cap B = \emptyset$.
- Let N denote the set of positive integers. Often, we will be concerned with a finitely additive, probability measure on the power set of N, P(N).
- μ is strongly continuous if for every ε > 0, there exists a measurable partition {F₁,..., F_n} of T such that μ(F_i) < ε for every i.
- If μ is strongly continuous then it is atomless. A countably additive measure μ is strongly continuous iff it is atomless.
- ► The range of a strongly continuous measure is convex.

Games on Finitely Additive Spaces

- ► Let $E = \{e^1, \dots, e^L\}$ be the set of unit vectors in \mathbb{R}^L and $S = \{s \in \mathbb{R}_+^L : \sum_{k=1}^L s_k = 1\}$ the unit simplex in \mathbb{R}^L .
- Let U be the set of real valued continuous functions defined on E × S, endowed with sup norm.
- Let (T, T, μ) be a finitely additive probability space.
- A *game* is a measurable function $\mathcal{G}: \mathcal{T} \longrightarrow \mathcal{U}.$
- A pure strategy profile (of a game \mathcal{G}) is a measurable function $f: T \longrightarrow E$.
- A f : T → E is a (pure strategy) Nash equilibrium of G if for almost all t,

$$\mathcal{G}(t)\left(f(t),\int_{\mathcal{T}}f\;\mathsf{d}\mu\right)\geq\mathcal{G}(t)\left(\mathsf{a},\int_{\mathcal{T}}f\;\mathsf{d}\mu\right)\;\text{for all}\;\mathsf{a}\in\mathsf{E}.$$

Mixed Strategies

- Pure strategy profile: $f : T \longrightarrow E$.
- Mixed strategy profile: $f : T \longrightarrow S$.
- ► Given a mixed strategy profile f and y ∈ S, the payoff to player t is

$$\mathcal{G}(t)\left(y,\int_{\mathcal{T}}f\,\mathrm{d}\mu\right)=\sum_{k=1}^{L}y_{k}\mathcal{G}(t)\left(e^{k},\int_{\mathcal{T}}f\,\mathrm{d}\mu\right).$$

A f : T → E is a (mixed strategy) Nash equilibrium of G if for almost all t,

$$\mathcal{G}(t)\left(f(t),\int_{\mathcal{T}}f\;\mathsf{d}\mu\right)\geq\mathcal{G}(t)\left(y,\int_{\mathcal{T}}f\;\mathsf{d}\mu\right)\;\text{for all }y\in\mathcal{S}.$$

Nonexistence of Nash Equilibria: Example

▶ Let $A = \{0, 1\}$ and K = [0, 1]. For each $t \in \mathbb{N}$, let the payoff function (on $A \times K$) be

$$\mathcal{G}(t)(a,x) = \left(x - rac{1}{t+1}
ight)^{a+1}, \ a \in A.$$

- We will derive the best responses and show that this game has no Nash equilibrium.
- Best responses:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/(t+1) \\ 1 & \text{if } x < 1/(t+1) \\ 0 & \text{if } x > 1/(t+1). \end{cases}$$

$$* x = 1/(t+1): \ \mathcal{G}(t)(0, x) = \mathcal{G}(t)(1, x) = 0.$$

$$* x < 1/(t+1): \ \mathcal{G}(t)(0, x) < 0 < \mathcal{G}(t)(1, x).$$

$$* x > 1/(t+1): \ 0 < \mathcal{G}(t)(0, x) < 1, \ \mathcal{G}(t)(1, x) = [\mathcal{G}(t)(0, x)]^2$$

Best responses:

$$\operatorname{argmax}_{a \in A} u_t(a, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/(t+1) \\ 1 & \text{if } x < 1/(t+1) \\ 0 & \text{if } x > 1/(t+1). \end{cases}$$

- Suppose that f from N to K is a Nash equilibrium. Let x = ∫_N f dµ.
 - ▶ If x = 0 then x < 1/(t+1) for all $t \in \mathbb{N}$ which implies that f(t) = 1 for all t and $\int_{\mathbb{N}} f d\mu = 1$, a contradiction.
 - If x > 0 then x > 1/(t+1) for almost all t (since the measure of a finite set is zero), which implies that f(t) = 0 for almost all t and $\int_{\mathbb{N}} f d\mu = 0$, again a contradiction.
- The game does not have a Nash equilibrium in pure or mixed strategies.

Nonexistence of Equilibria on General Measure Spaces

- Let T be a nonempty set and T a field of subsets of T.
 Let μ be a finitely additive probability measure on T.
- Assume that µ is not countably additive. We will show that there is a game on µ which has no pure or mixed strategy Nash equilibrium.

Claim

The following conditions are equivalent.

(i) μ is countably additive.

(ii) $\lim_{n\to\infty} \mu(A_n) = \mu(A)$ whenever $\{A_n\}$ is an increasing sequence of sets in \mathcal{T} with $A = \bigcup_{n=1}^{\infty} A_n \in \mathcal{T}$.

- Let $A = \{0, 1\}$ be the set of actions.
- Since µ is not countably additive, there is an increasing sequence of sets {B_n} in T such that

$$\cup_{n=1}^{\infty}B_n=T ext{ and } \lim_{n o\infty}\mu(B_n)=c<1.$$

- ▶ For $n \in \mathbb{N}$, let $C_1 = B_1$ and for $n \ge 2$, $C_n = B_n \setminus B_{n-1}$.
- ▶ $\{C_n\}$ is a sequence of pairwise disjoint sets and $\bigcup_{n=1}^{\infty} C_n = T$.
- Now we will define the payoffs. Let x ∈ [0, 1]. For each t ∈ C_n, let

$$\mathcal{G}(t)(a,x)=(x-\ell_n)^{a+1},\;a\in A$$
 where $\ell_n=c+rac{1-c}{n}.$

Note that ℓ₁ = 1, ℓ_n > c for each n and {ℓ_n} is a monotonically decreasing sequence converging to c.

The Example, contd.

• $\mathcal{G}(t)(a,x) = (x - \ell_n)^{a+1}$. Best responses:

$$\operatorname{argmax}_{a \in \mathcal{A}} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = \ell_n \\ 1 & \text{if } x < \ell_n \\ 0 & \text{if } x > \ell_n. \end{cases}$$

- ▶ Let $f : T \longrightarrow [0, 1]$ be a mixed strategy Nash equilibrium and $x = \int_T f \, d\mu$.
 - Suppose that x ≤ c < 1. Then for all t ∈ T, f(t) = 1 which implies that x = 1, a contradiction.</p>
 - ▶ Now suppose that x > c. Then there exists a unique $n_0 \in \mathbb{N}$ such that $\ell_{n_0+1} < x \leq \ell_{n_0}$. If $n \geq n_0 + 1$ and $t \in C_n$ then f(t) = 0. So, $x = \int_T f d\mu \leq \sum_{i=1}^{n_0} \mu(C_i) = \mu(B_{n_0}) \leq c$, a contradiction.
- The game does not have a Nash equilibrium in pure or mixed strategies.

A Sequence of Games

- Below we construct a sequence of games, each of which has a pure strategy Nash equilibrium.
- Fix any $n \in \mathbb{N}$, and let

$$\mathcal{G}^n(t) = egin{cases} \mathcal{G}(t) & ext{if } t \in B_n \ ar{u} & ext{if } t
ot \in B_n, \end{cases}$$

where $\bar{u}(0,x) = \bar{u}(1,x) = 0$ for any $x \in [0,1]$.

Let fⁿ be a function from T to A such that

$$f^n(t) = 1$$
 on B_n and $f^n(t) = 0$ on B_n^c .

- Then $\int_T f^n d\mu = \mu(B_n) \le c < \ell_n$ for every $n \in \mathbb{N}$.
- ► It is clear that $f^n(t)$ is a best response to $\int_T f^n d\mu$ for every player t,
- This means f^n is a Nash equilibrium of \mathcal{G}^n .

An Implication of Countable Additivity

Theorem

Let (T, T, μ) be a finitely additive probability space where T is a σ -algebra and μ is strongly continuous. Then the following are equivalent.

- Every game G on T with at least two actions has a mixed strategy Nash equilibrium.
- (ii) μ is countably additive.

Proof (*ii*) \Rightarrow (*i*). Schmeidler (1973).

 $(i) \Rightarrow (ii)$. Assume that μ is not countably additive. Consider the game \mathcal{G} in the example on general measure spaces. It does not have a pure/mixed strategy Nash equilibrium.

- The Nash equilibrium correspondence assigns the set of Nash equilibria to a game.
- Let Gⁿ, n ∈ N, and G be games on (T, T, μ). The Nash equilibrium correspondence has the *closed-graph property* if the following holds: if
 - $\{\mathcal{G}^n\}$ converges to \mathcal{G} pointwise,
 - f^n is a Nash equilibrium of \mathcal{G}^n for each n and
 - $\{f^n\}$ converges to f pointwise.

then f is a Nash equilibrium of G.

Another Implication of Countable Additivity

Theorem

Let (T, T, μ) be a finitely additive probability space where T is a σ -algebra and μ is strongly continuous. Then the following are equivalent.

- (i) The Nash equilibrium correspondence of games with at least two actions has the closed graph property.
- (ii) μ is countably additive.

Proof (*ii*) \Rightarrow (*i*). This is shown in Theorem 2 of Qiao-Yu-Zhang (2015).

 $(i) \Rightarrow (ii)$. Assume that μ is not countably additive. Consider the sequence of games $\{\mathcal{G}^n\}$ and the game \mathcal{G} in the example on general measure spaces. The closed graph property fails at \mathcal{G} .

epsilon-Equilibrium

- Let \mathcal{G} be a game on $(\mathcal{T}, \mathcal{T}, \mu)$ and $\epsilon > 0$.
- A strategy profile $f : T \longrightarrow S$ is an ϵ -equilibrium of \mathcal{G} if
 - there exists $T_{\epsilon} \in \mathcal{T}$ such that $\mu(T_{\epsilon}) \leq \epsilon$ and
 - for any $t \in T_{\epsilon}^{c}$,

$$\mathcal{G}(t)\left(f(t),\int_{\mathcal{T}}f\,\mathrm{d}\mu\right)\geq\mathcal{G}(t)\left(\mathsf{a},\int_{\mathcal{T}}f\,\mathrm{d}\mu\right)-\epsilon$$

for all $a \in E$.

Nonexistence of epsilon-Equilibrium

- The game is on \mathbb{N} , with $A = \{0, 1\}$ and K = [0, 1].
- For each player $t \in \mathbb{N}$, the payoff function is

 $\mathcal{G}(t)(0,x) = 0$ and $\mathcal{G}(t)(1,x) = 1 - 2^{t}x + 2^{t-1}$.

The best responses are:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = (1/2) + 2^{-t} \\ 1 & \text{if } x < (1/2) + 2^{-t} \\ 0 & \text{if } x > (1/2) + 2^{-t}. \end{cases}$$

► This game does not have an
e-equilibrium if 0 <
e ≤ 1/4.</p>

Existence of epsilon-Equilibrium

▶ A game \mathcal{G} on (T, \mathcal{T}, μ) is said to be *tight* if for any $\epsilon > 0$, there exist $\overline{T} \subseteq T$ such that $\mu(\overline{T}) < \epsilon$ and $\mathcal{G}(T \setminus \overline{T})$ is a relatively compact subset of \mathcal{U} .

Theorem

Let \mathcal{G} be a game on (T, \mathcal{T}, μ) . If the measure μ is strongly continuous and \mathcal{G} is tight, then \mathcal{G} has a pure strategy ϵ -equilibrium for every $\epsilon > 0$.