

Modeling Infinitely Many Agents: Why Countable Additivity Is Necessary

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Background and Motivation

- ▶ **Finite agent economies and games:** Arrow-Debreu (1954), McKenzie (1954), Nash (1951).
- ▶ **Economies and games with a continuum of agents:** Aumann (1964, 1966), Vind (1964), Schmeidler (1969, 1973).
- ▶ **Modeling individual negligibility:**
 - ▶ **Replication/Large finite approximations:** Edgeworth (1881), Debreu-Scarf (1963), Anderson (1978).
 - ▶ **Continuum models with an atomless measure:** Milnor-Shapley (1961), Aumann (1964), Schmeidler (1973), Hildenbrand (1974), Khan-Sun (2002).
 - ▶ **Infinitesimals, Loeb spaces:** Brown-Robinson (1972, 1975), Khan (1974), Brown-Loeb (1976), Khan-Sun (1996, 1999).
 - ▶ **Finitely additive economies:** Armstrong-Richter (1984, 1986), Weiss (1981), Feldman-Gilles (1985), Basile (1993).

Mathematical Preliminaries

- ▶ Let T be a nonempty set and \mathcal{T} a σ -algebra of subsets of T ,
 - (i) $T \in \mathcal{T}$,
 - (ii) $A \in \mathcal{T}$ implies $A^c \in \mathcal{T}$,
 - (iii) $A_n \in \mathcal{T}$ ($n = 1, 2, \dots$) implies $\cup_{n=1}^{\infty} A_n \in \mathcal{T}$.
- ▶ Let μ be a set function from \mathcal{T} to $[0, 1]$ with $\mu(T) = 1$.
 - ▶ μ is a **finitely additive measure** on \mathcal{T} if for any $A, B \in \mathcal{T}$ with $A \cap B = \emptyset$, $\mu(A \cup B) = \mu(A) + \mu(B)$.
 - ▶ μ is a **countably additive measure** on \mathcal{T} if for any sequence $\{A_n\}$ of pairwise disjoint sets in \mathcal{T} , $\mu(\cup_{n=1}^{\infty} A_n) = \sum_{n=1}^{\infty} \mu(A_n)$.
- ▶ The triple (T, \mathcal{T}, μ) will be called a (**finitely additive/countably additive**) **measure space**.

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- ▶ The triple (T, \mathcal{T}, μ) will be called a (**finitely additive/countably additive**) **measure space**.
- ▶ A measure μ is **atomless** if for every $\epsilon > 0$, there exists a \mathcal{T} -measurable partition $\{F_1, \dots, F_n\}$ of T such that $\mu(F_i) < \epsilon$ for every i .
- ▶ Let \mathbb{N} be the set of positive integers and $\mathcal{P}(\mathbb{N})$ its power set. **There are finitely additive, atomless measures on $\mathcal{P}(\mathbb{N})$** (such as a density charge).

Preview of the Results

- ▶ **Negative results on finitely additive spaces.**
 - ▶ An economy may not have a competitive equilibrium.
(Two examples)
 - ▶ A game may not have a Nash equilibrium.
(Two examples)
 - ▶ An economy may not have the idealized limit property.
 - ▶ A game may not have the idealized limit property.
- ▶ **Consequences.**
 - ▶ Necessity of countably additivity for economies:
both existence and idealized limit property hold.
 - ▶ Necessity of countably additivity for games:
both existence and idealized limit property hold.
- ▶ **Approximate equilibria on finitely additive spaces.**
 - ▶ An economy may not have an approximate competitive equilibrium.
A tightness assumption is sufficient for existence.
 - ▶ A game may not have an approximate Nash equilibrium.
A tightness assumption is sufficient for existence.

Upper Hemicontinuity of the Integral

- ▶ Let (T, \mathcal{T}, μ) be an atomless, **countably additive measure** space and X a metric space.
- ▶ Let $F : T \times X \mapsto \mathbb{R}^n$ be a correspondence.
- ▶ If $F(\cdot, x)$ is **measurable** and $F(t, \cdot)$ is **upper hemicontinuous** then

$$\int_T F(\cdot, x) d\mu$$

is **upper hemicontinuous** (in x).

- ▶ This results fails if μ is a finitely additive measure.

Lack of UHC under Integration

- ▶ Let $A = \{0, 1\}$ and $K = [0, 1]$.
Let μ be an atomless, **finitely additive measure** on $\mathcal{P}(\mathbb{N})$.
- ▶ Define a correspondence $F : \mathbb{N} \times K \mapsto A$ as:

$$F(t, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/t \\ 1 & \text{if } x < 1/t \\ 0 & \text{if } x > 1/t. \end{cases}$$

- ▶ Then

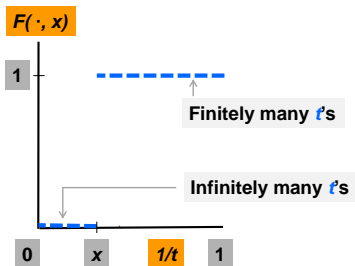
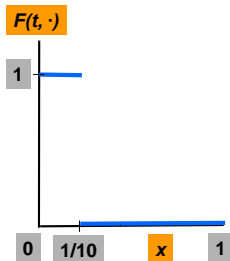
$$\int_{\mathbb{N}} F(\cdot, x) d\mu = \begin{cases} 1 & \text{if } x = 0 \\ 0 & \text{if } x > 0. \end{cases}$$

- ▶ Clearly, $\int_{\mathbb{N}} F(\cdot, x) d\mu$ is not uhc at $x = 0$.

Graphs of the Correspondence

$$F(t, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/t \\ 1 & \text{if } x < 1/t \\ 0 & \text{if } x > 1/t. \end{cases}$$

Let $t = 10$.



Example, contd.

- ▶ $F : \mathbb{N} \times K \mapsto A$.

$$F(t, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/t \\ 1 & \text{if } x < 1/t \\ 0 & \text{if } x > 1/t. \end{cases}$$

- ▶ Then

$$\int_{\mathbb{N}} F(\cdot, x) d\mu = \begin{cases} 1 & \text{if } x = 0 \\ 0 & \text{if } x > 0. \end{cases}$$

- ▶ Clearly, $\int_{\mathbb{N}} F(\cdot, x) d\mu$ is not uhc at $x = 0$.
- ▶ Let f be a measurable selection of $F(\cdot, x)$.
 - ▶ If $x = 0$ then $x < 1/t$ for all $t \in \mathbb{N}$, which implies that $f(t) = 1$ for all $t \in \mathbb{N}$ and $\int_{\mathbb{N}} f d\mu = 1$.
 - ▶ If $x > 0$ then $x > 1/t$ for almost all t , i.e., $f(t) = 0$ for almost all t and $\int_{\mathbb{N}} f d\mu = 0$.

Economies and Competitive Equilibria

- ▶ There are L goods and the commodity space is \mathbb{R}_+^L .
- ▶ Let \mathcal{U} denote the class of real valued, continuous utility functions on \mathbb{R}_+^L (endowed with the compact open topology).
- ▶ A $u \in \mathcal{U}$ is *strongly monotone* if $x \geq y$, $x \neq y$ implies that $u(x) > u(y)$.
- ▶ Let (T, \mathcal{T}, μ) be a finitely additive measure space. (space of agents)

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- ▶ Let (T, \mathcal{T}, μ) be a finitely additive measure space. (space of agents)
- ▶ An *economy* is a measurable mapping $\mathcal{E} = (u, \omega) : T \rightarrow \mathcal{U} \times \mathbb{R}_+^L$ such that ω is integrable and $\bar{\omega} = \int_T \omega \, d\mu \gg 0$.
- ▶ An *allocation* of \mathcal{E} is an integrable mapping f from T to \mathbb{R}_+^L .
An allocation is *feasible* if $\int_T f \, d\mu = \int_T \omega \, d\mu$.
- ▶ Given a price vector $p \in \mathbb{R}_+^L$, the *budget set* of consumer t is
$$B_t(p) = \{x \in \mathbb{R}_+^L : p \cdot x \leq p \cdot \omega_t\}.$$
- ▶ A *competitive equilibrium* of \mathcal{E} is a pair (p, f) , where $p \in \mathbb{R}_+^L \setminus \{0\}$, f is a feasible allocation and μ -a.e.;
 - (a) $f(t) \in B_t(p)$ and
 - (b) $u_t(f(t)) \geq u_t(x)$ for all $x \in B_t(p)$.
- ▶ An allocation f of \mathcal{E} is a *competitive allocation* if for some p , (p, f) is a competitive equilibrium.

Nonexistence of a CE: An Example on Integers

- ▶ The measure space is $(\mathbb{N}, \mathcal{P}(\mathbb{N}), \mu)$.
- ▶ The economy \mathcal{E} is defined as follows. For each $t \in \mathbb{N}$,

$$u_t(x_1, x_2) = \frac{t+1}{t} x_1^{\frac{t}{t+1}} + x_2, \quad \omega_t = \left(\frac{c+1}{2}, \frac{c+1}{2} \right),$$

where $0 \leq c < 1$.

- ▶ Assume that $(p_1, p_2) \in \mathbb{R}_+^2$ is a pair of equilibrium prices. We must have $p_1 > 0$ and $p_2 > 0$ since u_t is **strongly monotone** for each $t \in \mathbb{N}$.
- ▶ Without loss of generality suppose that $p_1 + p_2 = 1$.
- ▶ For any $t \in \mathbb{N}$, the unique solution of agent t 's problem;
maximize $u_t(x_1, x_2)$ **subject to** $p_1 x_1 + p_2 x_2 = (c+1)/2$ is

$$D_{t1} = \min \left\{ \frac{p_2^{t+1}}{p_1^{t+1}}, \frac{c+1}{2p_1} \right\}, \quad D_{t2} = \frac{c+1}{2p_2} - \frac{p_1 D_{t1}}{p_2}.$$

Nonexistence of a CE, continued

- ▶ $\omega_t = ((c + 1)/2, (c + 1)/2)$, $0 \leq c < 1$ for $t \in \mathbb{N}$.
- ▶ $p = (p_1, p_2) \gg 0$, $p_1 + p_2 = 1$. Income: $(c + 1)/2$.
- ▶ Demand functions:

$$D_{t1} = \min \left\{ \frac{p_2^{t+1}}{p_1^{t+1}}, \frac{c + 1}{2p_1} \right\}, \quad D_{t2} = \frac{c + 1}{2p_2} - \frac{p_1 D_{t1}}{p_2}.$$

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- ▶ **Case 1.** $p_2/p_1 < 1$. $\lim_{t \rightarrow \infty} D_{t1} = 0$. $\int_{\mathbb{N}} D_{t1} \, d\mu = 0$.

$$\int_{\mathbb{N}} D_{t2} \, d\mu = \frac{c + 1}{2p_2} > \frac{c + 1}{2} = \int_{\mathbb{N}} \omega_{t2} \, d\mu. \quad (\text{contradiction})$$

- ▶ **Case 2.** $p_2/p_1 \geq 1$. $p_2^{t+1}/p_1^{t+1} \geq 1$. $(c + 1)/(2p_1) \geq c + 1$.

Therefore, $D_{t1} \geq \min \{1, c + 1\} = 1$.

$$\int_{\mathbb{N}} D_{t1} \, d\mu \geq 1 > \frac{c + 1}{2} = \int_{\mathbb{N}} \omega_{t1} \, d\mu. \quad (\text{contradiction})$$

Nonexistence of a CE on General Measure Spaces

Claim

Let (T, \mathcal{T}, μ) be an atomless finitely additive measure space. Assume that μ is not countably additive. *Then there is an economy on (T, \mathcal{T}, μ) which has no competitive equilibrium.*

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Let (T, \mathcal{T}, μ) be an atomless finitely additive measure space. Assume that μ is not countably additive. Then there is an economy on (T, \mathcal{T}, μ) which has no competitive equilibrium.

- ▶ **Fact:** Let (T, \mathcal{T}, μ) be a finitely additive probability space. Then the following are equivalent.
 - μ is not countably additive.
 - There is an increasing sequence of sets $\{B_n\}$ in \mathcal{T} such that $\bigcup_{n=1}^{\infty} B_n = T$ and $\lim_{n \rightarrow \infty} \mu(B_n) = c < 1$.
- ▶ Since μ is not countably additive, there is an increasing sequence of sets $\{B_n\}$ in \mathcal{T} such that $\bigcup_{n=1}^{\infty} B_n = T$ and $\lim_{n \rightarrow \infty} \mu(B_n) = c < 1$.
- ▶ For $n \in \mathbb{N}$, let $C_1 = B_1$ and for $n \geq 2$, $C_n = B_n \setminus B_{n-1}$.
- ▶ $\{C_n\}$ is a sequence of pairwise disjoint sets and $\bigcup_{n=1}^{\infty} C_n = T$.

Nonexistence on General Measure Spaces, contd.

- Preferences and endowments: Let $t \in C_n$.

$$u_t(x_1, x_2) = \frac{n+1}{n} x_1^{\frac{n}{n+1}} + x_2, \quad \omega_t = \left(\frac{c+1}{2}, \frac{c+1}{2} \right).$$

- Assume that $p_1, p_2 \in \mathbb{R}_+^2$ is a pair of competitive equilibrium prices.
We must have $p_1 > 0$ and $p_2 > 0$ since u_t is **strongly monotone** for each t .
- Let $p_2 = 1$ and $p_1 > 0$. If $t \in C_n$, then

$$D_{t1} = \min \left\{ \frac{1}{p_1^{n+1}}, \frac{c+1}{2} \left(1 + \frac{1}{p_1} \right) \right\}, \quad D_{t2} = \frac{c+1}{2} (1 + p_1) - p_1 D_{t1}.$$

- To show that there is **no** competitive equilibrium, we will consider two cases: (i) $1 \geq p_1$ and (ii) $1 < p_1$.

Nonexistence of Equilibrium (Case 1)

Case 1: $1 \geq p_1$.

Let $t \in C_n$.

$$\begin{aligned} D_{t1} &= \min \left\{ \frac{1}{p_1^{n+1}}, \frac{c+1}{2} \left(1 + \frac{1}{p_1} \right) \right\} \geq \min \left\{ 1, \frac{c+1}{2} \left(1 + \frac{1}{p_1} \right) \right\} \\ &= \frac{c+1}{2} + \min \left\{ 1 - \frac{c+1}{2}, \frac{c+1}{2p_1} \right\}. \end{aligned}$$

Let $\theta = \min \left\{ 1 - \frac{c+1}{2}, \frac{c+1}{2p_1} \right\} > 0$.

Then $D_{t1} \geq \frac{c+1}{2} + \theta$ for any $t \in T$. Therefore,

$$\int_T D_{t1} \, d\mu \geq \frac{c+1}{2} + \theta > \frac{c+1}{2} = \int_T \omega_{t1} \, d\mu,$$

a contradiction.

Nonexistence of Equilibrium (Case 2)

Case 2: $1 < p_1$. Note that $D_{t1} \leq 1/p_1^{n+1}$ for any $t \in C_n$. Then

(i) $D_{t1} \leq 1$ for any $t \in T$ and

(ii) if $t \in C_{n+1}$ and $n > m$ then $D_{t1} \leq 1/p_1^{n+2} \leq 1/p_1^{m+2}$.

Fix a positive integer m .

$$\begin{aligned}\int_T D_{t1} \, d\mu &= \int_{B_m} D_{t1} \, d\mu + \int_{T \setminus B_m} D_{t1} \, d\mu \\ &\leq \int_{B_m} 1 \, d\mu + \int_{T \setminus B_m} \frac{1}{p_1^{m+2}} \, d\mu = \mu(B_m) + \frac{1}{p_1^{m+2}} \mu(T \setminus B_m).\end{aligned}$$

Observe that $\mu(T \setminus B_m) \geq 1 - c$ for any m .

Let m tend to infinity.

$$\int_T D_{t1} \, d\mu \leq \lim_{m \rightarrow \infty} \mu(B_m) = c.$$

$$\begin{aligned}\int_T D_{t2} \, d\mu &= \frac{c+1}{2}(1+p_1) - p_1 \int_T D_{t1} \, d\mu \\ &\geq \frac{c+1}{2}(1+p_1) - p_1 c > \frac{c+1}{2} = \int_T \omega_{t2} \, d\mu,\end{aligned}$$

a contradiction.

Games and Nash Equilibria

- ▶ Let $E = \{e^1, \dots, e^L\}$ be the set of unit vectors in \mathbb{R}^L and $S = \{s \in \mathbb{R}_+^L : \sum_{k=1}^L s_k = 1\}$ the unit simplex in \mathbb{R}^L .
- ▶ Let \mathcal{V} be the set of real valued continuous functions defined on $E \times S$, endowed with sup norm.
- ▶ (T, \mathcal{T}, μ) is an atomless, **countably/finitely additive probability space**.
- ▶ A *game* is a measurable function $\mathcal{G} : T \rightarrow \mathcal{V}$.
- ▶ A *pure strategy profile* is a measurable function $f : T \rightarrow E$.
- ▶ A $f : T \rightarrow E$ is a *pure strategy Nash equilibrium* of \mathcal{G} if μ -a.e.;

$$\mathcal{G}(t)(f(t), \int_T f \, d\mu) \geq \mathcal{G}(t)(a, \int_T f \, d\mu) \text{ for all } a \in E.$$

Games and Nash Equilibria, contd.

- ▶ **Pure strategy profile:** $f : T \rightarrow E$.
- ▶ **Mixed strategy profile:** $g : T \rightarrow S$.
- ▶ Given a mixed strategy profile g and $y \in S$, the *payoff to player t* is

$$\mathcal{G}(t) \left(y, \int_T g \, d\mu \right) = \sum_{k=1}^L y_k \mathcal{G}(t) \left(e^k, \int_T g \, d\mu \right).$$

- ▶ A $g : T \rightarrow S$ is a *mixed strategy Nash equilibrium* of \mathcal{G} if μ -a.e.;

$$\mathcal{G}(t) \left(g(t), \int_T g \, d\mu \right) \geq \mathcal{G}(t) \left(y, \int_T g \, d\mu \right) \text{ for all } y \in S.$$

Existence of Nash Equilibrium

Theorem (Schmeidler)

Every finite action game on a countably additive measure space has a pure strategy Nash equilibrium.

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Theorem (Schmeidler)

Every *finite action game on a countably additive measure space has a pure strategy Nash equilibrium.*

- ▶ Define a correspondence $B : T \times S \mapsto E$ by

$$B(t, s) = \{e^k \in E : \mathcal{G}(t)(e^k, s) \geq \mathcal{G}(t)(a, s) \text{ for all } a \in E\}.$$

- ▶ $B(t, s)$ is nonempty, $B(\cdot, s)$ is measurable and $B(t, \cdot)$ is uhc.
- ▶ Let $\Gamma(s) = \int_T B(\cdot, s) d\mu$.
 - ▶ $\Gamma(s)$ is nonempty for each $s \in S$.
 - ▶ $\Gamma(\cdot)$ is uhc (integration preserves uhc).
 - ▶ $\Gamma(\cdot)$ is convex valued (by Lyapunov's theorem).
- ▶ Γ has a fixed point s^* (by Kakutani's fixed point theorem).
- ▶ There is $f : T \rightarrow E$ such that $\int_T f d\mu = s^*$ and μ -a.e., $f(t) \in B(t, s^*)$.
- ▶ This f is a pure strategy Nash equilibrium of \mathcal{G} . ■

Nonexistence of an NE: An Example on Integers

- ▶ Let $A = \{0, 1\}$ and $K = [0, 1]$.
Any $x \in K$ can be interpreted as the weight on action 1.
- ▶ For each $t \in \mathbb{N}$, let the payoff function on $A \times K$ is

$$\mathcal{G}(t)(a, x) = a \left(\frac{1}{t} - x \right), \quad a \in A.$$

- ▶ We will show that this game has no Nash equilibrium.
- ▶ Best responses:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/t \\ 1 & \text{if } x < 1/t \\ 0 & \text{if } x > 1/t. \end{cases}$$

- ▶ $x = 1/t$: $\mathcal{G}(t)(0, x) = \mathcal{G}(t)(1, x) = 0$.
- ▶ $x < 1/t$: $\mathcal{G}(t)(0, x) = 0 < \mathcal{G}(t)(1, x)$.
- ▶ $x > 1/t$: $\mathcal{G}(t)(0, x) = 0 > \mathcal{G}(t)(1, x)$.

Example, contd.

- ▶ Best responses:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = 1/t \\ 1 & \text{if } x < 1/t \\ 0 & \text{if } x > 1/t. \end{cases}$$

- ▶ Suppose that f from \mathbb{N} to K is a (mixed) Nash equilibrium.

Let $x = \int_{\mathbb{N}} f \, d\mu$.

- ▶ If $x = 0$ then $x < 1/t$ for all $t \in \mathbb{N}$ which implies that $f(t) = 1$ for all t and $\int_{\mathbb{N}} f \, d\mu = 1$. (contradiction)
- ▶ If $x > 0$ then $x > 1/t$ for almost all t (since the measure of a finite set is zero), which implies that $f(t) = 0$ for almost all t and $\int_{\mathbb{N}} f \, d\mu = 0$. (contradiction)
- ▶ The game does not have a Nash equilibrium in pure or mixed strategies.

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Claim

Let (T, \mathcal{T}, μ) be an atomless finitely additive measure space. Assume that μ is not countably additive. Then there is a game on (T, \mathcal{T}, μ) which has no Nash equilibrium.

- ▶ Since μ is not countably additive, there is an increasing sequence of sets $\{B_n\}$ in \mathcal{T} such that

$$\bigcup_{n=1}^{\infty} B_n = T \text{ and } \lim_{n \rightarrow \infty} \mu(B_n) = c < 1.$$

- ▶ For $n \in \mathbb{N}$, let $C_1 = B_1$ and for $n \geq 2$, $C_n = B_n \setminus B_{n-1}$.
- ▶ $\{C_n\}$ is a sequence of pairwise disjoint sets and $\bigcup_{n=1}^{\infty} C_n = T$.
- ▶ $A = \{0, 1\}$, $K = [0, 1]$. For each $t \in C_n$, let

$$\mathcal{G}(t)(a, x) = a(\ell_n - x), \quad \text{where} \quad \ell_n = c + \frac{1-c}{n}.$$

- ▶ Note that $\ell_1 = 1$, $\ell_n > c$ for each n and $\{\ell_n\}$ is a monotonically decreasing sequence converging to c .

The Example, contd.

▶ $\mathcal{G}(t)(a, x) = a(\ell_n - x)$. $\ell_n = c + [(1 - c)/n]$.

- ▶ Best responses:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = \ell_n \\ 1 & \text{if } x < \ell_n \\ 0 & \text{if } x > \ell_n. \end{cases}$$

- ▶ Let $f : T \rightarrow [0, 1]$ be a mixed strategy Nash equilibrium and $x = \int_T f \, d\mu$.

- ▶ Suppose that $x \leq c < 1$. Then $x < \ell_n$ for all n .
For all $t \in T$, $f(t) = 1$, i.e., $x = 1$. (contradiction)

- ▶ Now suppose that $x > c$.

There is a unique $n_0 \in \mathbb{N}$ such that $\ell_{n_0+1} < x \leq \ell_{n_0}$.

If $n \geq n_0 + 1$ and $t \in C_n$ then $f(t) = 0$.

So, $x = \int_T f \, d\mu \leq \sum_{i=1}^{n_0} \mu(C_i) = \mu(B_{n_0}) \leq c$. (contradiction)

- ▶ The game does not have a Nash equilibrium in pure or mixed strategies.

Definition

A measurable mapping $\alpha^m : T \rightarrow \{1, \dots, m\}$ is a *replication function* if $\mu(\alpha^m)^{-1}(\{i\}) = 1/m$ for $i = 1, \dots, m$.

Definition

An economy \mathcal{E} on an atomless finitely additive measure space (T, \mathcal{T}, μ) is said to have the *idealized limit property* if

- (1) for any sequence $\{\mathcal{E}^n\}_{n=1}^{\infty}$ of finite-agent economies with $\{f_n\}_{n=1}^{\infty}$ as competitive allocations, where the number of agents in \mathcal{E}^n is k_n and $\lim_{n \rightarrow \infty} k_n = \infty$,
- (2) for any sequence of replication functions $\{\alpha^{k_n}\}_{n=1}^{\infty}$ such that $\mathcal{E}^n \circ \alpha^{k_n}$ converges to \mathcal{E} pointwise on T , $f^n \circ \alpha^{k_n}$ converges to some allocation f pointwise on T , and $\lim_{n \rightarrow \infty} \int_T \omega^n \circ \alpha^{k_n} d\mu = \int_T \omega d\mu$, then f is a competitive allocation of \mathcal{E} .

Example: No Idealized Limit

- ▶ Consider the economy $\mathcal{E} = (u, \omega)$, for each $t \in \mathbb{N}$,

$$u_t(x_1, x_2) = \frac{t+1}{t} x_1^{\frac{t}{t+1}} + x_2, \quad \omega_t = \left(\frac{c+1}{2}, \frac{c+1}{2} \right),$$

where $0 \leq c < 1$.

- ▶ Fix any $n \in \mathbb{N}$. Let \mathcal{E}^n be the restriction of \mathcal{E} on $\{1, \dots, n\}$.
- ▶ Since \mathcal{E}^n is a finite economy with concave and strictly increasing utility functions, there exists a competitive equilibrium f^n .

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- ▶ Fix any $n \in \mathbb{N}$. Let \mathcal{E}^n be the restriction of \mathcal{E} on $\{1, \dots, n\}$.
- ▶ Since \mathcal{E}^n is a finite economy with concave and strictly increasing utility functions, there exists a competitive equilibrium f^n .
- ▶ Let $\{A_k^n\}_{k=1}^n$ be a partition of \mathbb{N} such that $A_k^n = \{mn + k : m = 0, 1, \dots\}$.
- ▶ Let $\alpha^n(t) = k$ for any $t \in A_k^n$, where $k = 1, \dots, n$.
- ▶ Note that for any $n \geq t$, $t \in A_t^n$. Then $u_{\alpha^n(t)} = u_t$ for any $n \geq t$, which implies $\mathcal{E}^n \circ \alpha^n$ converges to \mathcal{E} pointwise.
- ▶ Moreover, $f^n \circ \alpha^n$ converges pointwise and $\lim_{n \rightarrow \infty} \int_{\mathcal{T}} \omega^n \circ \alpha^n d\mu = \int_{\mathcal{T}} \omega d\mu$.
- ▶ However, the limit economy $\mathcal{E} = (u, \omega)$ has no competitive equilibrium, which implies $\mathcal{E} = (u, \omega)$ does not have the idealized limit property.

Definition

A game \mathcal{G} on an atomless finitely additive measure space (T, \mathcal{T}, μ) is said to have the *idealized limit property* if

- (1) for any sequence $\{\mathcal{G}^n\}_{n=1}^{\infty}$ of finite-agent games with $\{f_n\}_{n=1}^{\infty}$ as pure strategy Nash equilibria, where the number of agents in \mathcal{G}^n is k_n and $\lim_{n \rightarrow \infty} k_n = \infty$,
- (2) for any sequence of replication functions $\{\alpha^{k_n}\}_{n=1}^{\infty}$ such that $\mathcal{G}^n \circ \alpha^{k_n}$ converges to \mathcal{G} pointwise on T , and $f^n \circ \alpha^{k_n}$ converges to some pure strategy profile f pointwise on T ,

then f is a pure strategy Nash equilibrium of \mathcal{G} .

The next example shows that the idealized limit property may fail for a game with countably many agents.

Example: No Idealized Limit

- ▶ Consider the game \mathcal{G} , for $t \in \mathbb{N}$, $\mathcal{G}(t)(a, x) = a[(1/t) - x]$
- ▶ Fix any $n \in \mathbb{N}$. Let \mathcal{G}^n be the restriction of \mathcal{G} on $\{1, \dots, n^2\}$.
- ▶ Let $\{A_i^n\}_{i=1}^{n^2}$ be a partition of \mathbb{N} such that
$$A_i^n = \{mn^2 + i : m = 0, 1, \dots\}.$$
- ▶ Let $\alpha^{n^2}(t) = k$ for any $t \in A_k^n$, where $k = 1, \dots, n^2$.
Note that for any $n \geq \sqrt{t}$, $\alpha^{n^2}(t) = t$.
- ▶ Then $\mathcal{G}^n \circ \alpha^{n^2}(t) = \mathcal{G}(t)$ for any $n \geq \sqrt{t}$, which implies $\mathcal{G}^n \circ \alpha^{n^2}$ converges to \mathcal{G} pointwise on T .

Example: No Idealized Limit

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- ▶ Then $\mathcal{G}^n \circ \alpha^{n^2}(t) = \mathcal{G}(t)$ for any $n \geq \sqrt{t}$, which implies $\mathcal{G}^n \circ \alpha^{n^2}$ converges to \mathcal{G} pointwise on T .
- ▶ Fix any $n \geq 2$. Let

$$f^n(i) = \begin{cases} 1 & \text{if } i \leq n \\ 0 & \text{if } i > n. \end{cases}$$

Then f^n is a Nash equilibrium of \mathcal{G}^n .

- ▶ Fix any $t \in \mathbb{N}$. For any $n \geq t$, $\alpha^{n^2}(t) = t$, which implies $f^n \circ \alpha^{n^2}(t) = 1$.
Then $f^n \circ \alpha^{n^2}(t) \rightarrow 1$ as $n \rightarrow \infty$.
- ▶ However, the limit game \mathcal{G} has no mixed strategy Nash equilibrium.

Necessity of Countable Additivity: Economies

We have seen failures of both existence and the idealized limit property for competitive equilibria in economies over a finitely additive measure space.

The next theorem shows the equivalence of countable additivity of the agent space with the validity of each of the properties.

Theorem

Let (T, \mathcal{T}, μ) be a finitely additive measure space. Assume that all the preferences are strongly monotone. Then the following statements hold.

- (i) Every economy \mathcal{E} on (T, \mathcal{T}, μ) has a competitive equilibrium if and only if μ is countably additive.*
- (ii) Every economy \mathcal{E} on (T, \mathcal{T}, μ) has the idealized limit property if and only if μ is countably additive.*

CA \Rightarrow Existence: Aumann (1966).

Existence \Rightarrow CA: Earlier example.

CA \Rightarrow ILP: Proof in the paper. (Follows Hildenbrand (1974))

ILP \Rightarrow CA: Earlier example on \mathbb{N} can be modified to any T .

Necessity of Countable Additivity: Games

We have seen failures of both existence and the idealized limit property for Nash equilibria in games over a finitely additive measure space.

The next theorem shows the equivalence of countable additivity of the agent space with the validity of each of the properties.

Theorem

Let (T, \mathcal{T}, μ) be a finitely additive measure space. Then the following statements hold.

- (i) Every game \mathcal{G} on (T, \mathcal{T}, μ) has a pure strategy Nash equilibrium if and only if μ is countably additive.
- (ii) Every game \mathcal{G} on (T, \mathcal{T}, μ) has idealized limit property if and only if μ is countably additive.

CA \Rightarrow Existence: Schmeidler (1973).

Existence \Rightarrow CA: Earlier example.

CA \Rightarrow ILP: Proof in the paper.

ILP \Rightarrow CA: Earlier example on \mathbb{N} can be modified to any T .

Approximate Competitive Equilibria

Earlier, we have seen examples that an economy may not have a competitive equilibrium. It is natural to ask if approximate competitive equilibria exist.

Definition

Let \mathcal{E} be an economy on (T, \mathcal{T}, μ) and $\epsilon > 0$. (p, f) is an ϵ -competitive equilibrium of \mathcal{E} if $p \in \mathbb{R}_+^L \setminus \{0\}$, f is a feasible allocation, $f(t) \in B_t(p)$ for almost all t and there exists $T_\epsilon \in \mathcal{T}$ such that:

- (a) $\mu(T_\epsilon) \leq \epsilon$ and
- (b) for almost all $t \in T_\epsilon^c$, $u_t(f(t)) \geq u_t(y) - \epsilon$ for any $y \in B_t(p)$.

In general, an ϵ -competitive equilibrium may not exist, as shown by the next Example.

Nonexistence of Approximate Competitive Equilibria

- ▶ The economy is on \mathbb{N} .
- ▶ The utility function and endowment of $t \in \mathbb{N}$ is,

$$u_t(x_1, x_2) = e^t \left[\frac{t+1}{t} x_1^{\frac{t}{t+1}} + x_2 \right], \quad \omega_t = \left(\frac{c+1}{2}, \frac{c+1}{2} \right),$$

where $0 \leq c < 1/3$.

- ▶ This economy does not have an ϵ -competitive equilibrium if $0 < \epsilon \leq 1/3$.

Existence of Approximate Competitive Equilibria

Definition

An economy \mathcal{E} on (T, \mathcal{T}, μ) is *tight* if for any $\epsilon > 0$, there exists $\bar{T} \subseteq T$ such that

- (a) $\mu(\bar{T}) < \epsilon$ and
- (b) $\mathcal{E}(T \setminus \bar{T})$ is a relatively compact subset of $\mathcal{U} \times \mathbb{R}^L_+$.

Proposition

If an economy \mathcal{E} is tight, then it has an ϵ -competitive equilibrium for every $\epsilon > 0$.

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Proposition

If an economy \mathcal{E} is tight, then it has an ϵ -competitive equilibrium for every $\epsilon > 0$.

- ▶ The existence of an ϵ -competitive equilibrium for every $\epsilon > 0$ does not imply that there is a competitive equilibrium. We demonstrate this by means of an earlier example.

Approximate Competitive Equilibria in an Example

- ▶ Take $c = 0$ in the first example. The (tight) economy is

$$u_t(x_1, x_2) = \frac{t+1}{t} x_1^{\frac{t}{t+1}} + x_2, \quad \omega_t = \left(\frac{1}{2}, \frac{1}{2} \right).$$

- ▶ If $p \gg 0$ and $p_1 + p_2 = 1$, then the demand functions are

$$D_{t1} = \min \left\{ \frac{p_2^{t+1}}{p_1^{t+1}}, \frac{1}{2p_1} \right\}, \quad D_{t2} = \frac{1}{2p_2} - \frac{p_1 D_{t1}}{p_2}.$$

- ▶ Let $p = (1/2, 1/2)$ and $f(t) = (1/2, 1/2) = \omega_t$.

For any $\epsilon > 0$, (p, f) is an ϵ -competitive equilibrium.

- ▶ $D_{t1} = 1$ and $D_{t2} = 0$. The maximized utility is $(t+1)/t$.

For each t , $f(t)$ is in the budget set and f is a feasible allocation.

- ▶ We will show that for any $\epsilon > 0$, and for almost all t ,

$$\frac{t+1}{t} \left(\frac{1}{2} \right)^{\frac{t}{t+1}} + \frac{1}{2} > \frac{t+1}{t} - \epsilon, \quad \epsilon > \frac{t+1}{t} - \frac{t+1}{t} \left(\frac{1}{2} \right)^{\frac{t}{t+1}} - \frac{1}{2}.$$

- ▶ As t tends to infinity, the RHS tends to zero. So, given $\epsilon > 0$, there exists $t_0 \in \mathbb{N}$ such that for all $t \geq t_0$, the above inequality holds.

Approximate Nash Equilibria

Earlier, we have seen examples that a game may not have a Nash equilibrium. It is natural to ask whether approximate Nash equilibria exist.

Definition

Let \mathcal{G} be a game on (T, \mathcal{T}, μ) and $\epsilon > 0$. A strategy profile $f : T \rightarrow S$ is an ϵ -Nash equilibrium of \mathcal{G} if there exists $T_\epsilon \in \mathcal{T}$ such that

- (a) $\mu(T_\epsilon) \leq \epsilon$ and
- (b) for almost all $t \in T_\epsilon^c$, $\mathcal{G}(t)(f(t), \int_T f d\mu) \geq \mathcal{G}(t)(y, \int_T f d\mu) - \epsilon$ for any $y \in S$.

In general, an ϵ -Nash equilibrium may not exist, as shown by the next Example.

Nonexistence of Approximate Nash Equilibria

- ▶ The game is on \mathbb{N} , with $A = \{0, 1\}$ and $K = [0, 1]$.
- ▶ For each player $t \in \mathbb{N}$, the payoff function is $\mathcal{G}(t)(0, x) = 0$ and

$$\mathcal{G}(t)(1, x) = \begin{cases} 1 + 2^{t-1}(1 - 2x) & \text{if } -1 \leq 1 + 2^{t-1}(1 - 2x) \leq 1 \\ 1 & \text{if } 1 + 2^{t-1}(1 - 2x) > 1 \\ -1 & \text{if } 1 + 2^{t-1}(1 - 2x) < -1. \end{cases}$$

- ▶ The **best responses** are:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = (1/2) + 2^{-t} \\ 1 & \text{if } x < (1/2) + 2^{-t} \\ 0 & \text{if } x > (1/2) + 2^{-t}. \end{cases}$$

- ▶ This game does not have an ϵ -Nash equilibrium if $0 < \epsilon \leq 1/4$.

Nonexistence of Approximate NE: Case 1

- ▶ Best responses:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = (1/2) + 2^{-t} \\ 1 & \text{if } x < (1/2) + 2^{-t} \\ 0 & \text{if } x > (1/2) + 2^{-t}. \end{cases}$$

- ▶ Let $0 < \epsilon \leq 1/4$ and suppose that f from \mathbb{N} to $[0, 1]$ is an ϵ -equilibrium. Then there exists $I_\epsilon \subseteq \mathbb{N}$ such that, $\mu(I_\epsilon) \leq \epsilon$ and for any $t \in I_\epsilon^c$,

$$\mathcal{G}(t)(f(t), x) \geq \max\{\mathcal{G}(t)(0, x), \mathcal{G}(t)(1, x)\} - \epsilon,$$

where $x = \int_T f \, d\mu$.

Nonexistence of Approximate NE: Case 1

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where $x = \int_T f \, d\mu$.

- ▶ **Case 1.** $x \leq 1/2$. For all $t \in \mathbb{N}$, $1 + 2^{t-1}(1 - 2x) \geq 1$.
 $\mathcal{G}(t)(1, x) = 1 > \mathcal{G}(t)(0, x)$. Therefore, for any $t \in I_\epsilon^c$,
 $\mathcal{G}(t)(f(t), x) \geq 1 - \epsilon$, which means $f(t) \geq 1 - \epsilon$.

$$x = \int_{I_\epsilon} f \, d\mu + \int_{I_\epsilon^c} f \, d\mu \geq \int_{I_\epsilon^c} f \, d\mu \geq (1 - \epsilon)^2 > \frac{1}{2}. \quad \text{contradiction}$$

Nonexistence of Approximate NE: Case 2

- ▶ Best responses:

$$\operatorname{argmax}_{a \in A} \mathcal{G}(t)(a, x) = \begin{cases} \{0, 1\} & \text{if } x = (1/2) + 2^{-t} \\ 1 & \text{if } x < (1/2) + 2^{-t} \\ 0 & \text{if } x > (1/2) + 2^{-t}. \end{cases}$$

- ▶ $0 < \epsilon \leq 1/4$. There is $I_\epsilon \subseteq \mathbb{N}$ such that, $\mu(I_\epsilon) \leq \epsilon$ and for any $t \in I_\epsilon^c$,

$$\mathcal{G}(t)(f(t), x) \geq \max\{\mathcal{G}(t)(0, x), \mathcal{G}(t)(1, x)\} - \epsilon,$$

where $x = \int_T f \, d\mu$.

Nonexistence of Approximate NE: Case 2

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- ▶ $0 < \epsilon \leq 1/4$. There is $I_\epsilon \subseteq \mathbb{N}$ such that, $\mu(I_\epsilon) \leq \epsilon$ and for any $t \in I_\epsilon^c$,

$$\mathcal{G}(t)(f(t), x) \geq \max\{\mathcal{G}(t)(0, x), \mathcal{G}(t)(1, x)\} - \epsilon,$$

where $x = \int_T f \, d\mu$.

- ▶ **Case 2.** $x > 1/2$. For almost all $t \in I_\epsilon^c$, $1 + 2^{t-1}(1 - 2x) < -1$.
 $\mathcal{G}(t)(0, x) > \mathcal{G}(t)(1, x) = -1$. Therefore, for all $t \in I_\epsilon^c$,
 $\mathcal{G}(t)(f(t), x) \geq -\epsilon$, which means $f(t) \leq \epsilon$.

$$x = \int_{I_\epsilon} f \, d\mu + \int_{I_\epsilon^c} f \, d\mu \leq \mu(I_\epsilon) + \epsilon(1 - \epsilon) \leq \epsilon + \epsilon \leq \frac{1}{2}. \quad \text{contradiction}$$

Existence of Approximate Nash Equilibria

Definition

A game \mathcal{G} on (T, \mathcal{T}, μ) is *tight* if for any $\epsilon > 0$, there exists $\bar{T} \subseteq T$ such that

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Proposition

If a game is \mathcal{G} is tight, then it has a pure strategy ϵ -Nash equilibrium for every $\epsilon > 0$.

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Proposition

If a game is \mathcal{G} is tight, then it has a pure strategy ϵ -Nash equilibrium for every $\epsilon > 0$.

- ▶ The existence of an ϵ -Nash equilibrium for every $\epsilon > 0$ does not ensure the existence of an NE. **Example:** $\mathcal{G}(t)(a, x) = a[(1/t) - x]$ on \mathbb{N} .
- ▶ The game is tight. It has an ϵ -Nash equilibrium for every $\epsilon > 0$.
- ▶ Explicitly, $f(t) = 0$ for all $t \in \mathbb{N}$ is an ϵ -Nash equilibrium.
 $\mathcal{G}(t)(0, 0) = 0$, $\mathcal{G}(t)(1, 0) = 1/t$, $0 \geq (1/t) - \epsilon$ for almost all t .
- ▶ However, as has been shown, the game does not have a Nash equilibrium.

Summary of Results

- ▶ **Negative results on finitely additive spaces.**
 - ▶ An economy may not have a competitive equilibrium.
(Two examples)
 - ▶ A game may not have a Nash equilibrium.
(Two examples)
 - ▶ An economy may not have the idealized limit property.
 - ▶ A game may not have the idealized limit property.
- ▶ **Consequences.**
 - ▶ Necessity of countably additivity for economies:
both existence and idealized limit property hold.
 - ▶ Necessity of countably additivity for games:
both existence and idealized limit property hold.
- ▶ **Approximate equilibria on finitely additive spaces.**
 - ▶ An economy may not have an approximate competitive equilibrium.
A tightness assumption is sufficient for existence.
 - ▶ A game may not have an approximate Nash equilibrium.
A tightness assumption is sufficient for existence.